

Lecture 40

TM Example, Encoding of a TM

Turing Machines: An Example

Example: Design a TM that decides $L = \{0^n 1^n 2^n \mid n \geq 1\}$ over $\{0,1,2\}$.

Solution: Idea: First ensure that input initially contains a few 0s, followed by a few 1s, followed by a few 2s. Then, check that the number of 0s, 1s, and 2s are equal.

$$Q = \{q_s, q_a, q_r, \dots\}$$

$$\Sigma = \{0,1,2\} \quad \Gamma = \{0,1,2, \sqcup, \dots\}$$

δ : Ensuring that input contains 0s followed by 1s followed by 2s.

$$\delta(q_s, 0) = (q_{find_1}, 0, R), \quad \delta(q_s, 1) = (q_r, 1, R), \quad \delta(q_s, 2) = (q_r, 2, R), \quad \delta(q_s, \sqcup) = (q_r, \sqcup, R)$$

$$\delta(q_{find_1}, 0) = (q_{find_1}, 0, R), \quad \delta(q_{find_1}, 1) = (q_{find_2}, 1, R), \quad \delta(q_{find_1}, 2) = (q_r, 2, R), \quad \delta(q_{find_1}, \sqcup) = (q_r, \sqcup, R)$$

$$\delta(q_{find_2}, 0) = (q_r, 0, R), \quad \delta(q_{find_2}, 1) = (q_{find_2}, 1, R), \quad \delta(q_{find_2}, 2) = (q_{find_\sqcup}, 2, R), \quad \delta(q_{find_2}, \sqcup) = (q_r, \sqcup, R)$$

$$\delta(q_{find_\sqcup}, 0) = (q_r, 0, R), \quad \delta(q_{find_\sqcup}, 1) = (q_r, 1, R), \quad \delta(q_{find_\sqcup}, 2) = (q_{find_\sqcup}, 2, R), \quad \delta(q_{find_\sqcup}, \sqcup) = (q_{cr_2}, \sqcup, L) \dots$$

Turing Machines: An Example

Example: Design a TM that decides $L = \{0^n 1^n 2^n \mid n \geq 1\}$

Solution: δ : Overview of checking that input contains equal number of 0s, 1s, and 2s.

- ▶ Replace the rightmost 2 by Z, rightmost 1 by Y, and rightmost 0 by X, in one iteration, repeatedly.
- ▶ Halt with q_r if machine can replace one out of $\{0,1,2\}$ but not the others.
- ▶ Halt with q_a if machine has replaced all of $\{0,1,2\}$.



Encoding a TM

Goal: Encode a TM M as a binary string x such that M can be recovered from x .

Warm Up: Encoding an integer:

$3 \rightarrow 11, 5 \rightarrow 101, 45 \rightarrow 101101$ (*Encode using binary representation*)

Encoding a pair (or tuple) of integers:

$\langle 5, 2 \rangle = 10110$ (*Concatenate the binary representation of 5 and 2*)

$\langle 5, 2 \rangle$ cannot be recovered from 10110 as 10110 is also the encoding of $\langle 2, 6 \rangle$.

Scheme: Encode $\langle x, y \rangle$ as $rep(x)01rep(y)$, where $rep(x)$, $rep(y)$ are the binary strings obtained from replacing every 0 with 00 and every 1 with 11 in the binary representation of x , y , respectively.

$\langle 5, 2 \rangle = 110011011100$ (*01 acts like a separator of two numbers*)

Encoding a TM

Observation: Encoding δ of Turing machine M is sufficient to encode M .

Let $f: Q \rightarrow \mathbb{Z}^+$, $g: \Gamma \rightarrow \mathbb{Z}^+$, $h: \{L, R\} \rightarrow \{1, 2\}$ be injective functions from Q , Γ , and $\{L, R\}$ to integers such that:

- ▶ f maps q_{start} , q_{accept} , q_{reject} to 1, 2, 3, respectively.
- ▶ g maps 0, 1, \sqcup to 1, 2, 3, respectively.

We can encode one entry of δ the following way:

$$\delta(q, a) = (q', b, M) \rightarrow \text{rep}(f(q))\mathbf{01}\text{rep}(g(a))\mathbf{01}\text{rep}(f(q'))\mathbf{01}\text{rep}(g(b))\mathbf{01}\text{rep}(h(M))$$

$$\text{E.g. } \delta(q_{start}, 0) = (q_{accept}, 1, R) \rightarrow 11\mathbf{01}11\mathbf{01}1100\mathbf{01}1100\mathbf{01}1100$$

Complete δ can be encoded by concatenating encoding of all entries using 01 as separator.

Note: Binary strings not in required format such as 1010 correspond to some trivial TM.